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# Zero-Knowledge Proofs for browser-based private authentication to blockchain

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# Using existing credentials and MPC-in-the-Head for browser-based private authentication to blockchain

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: Cybernetica AS



**ZKPROOF**

Paving the path to adoption

8th Workshop

Rome, Italy

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# Authentication in blockchain

## Problem

- User wants to post a transaction to the blockchain
  - Authenticated, i.e. signed
- User does not want to manage the private keys

## Solution

- Get a JWT from some OAuth provider (Google, Apple, ...)
  - Authenticates the user
  - Contains the transaction
- Post the JWT to blockchain

## New problem

JWT payload has sensitive fields (name, e-mail, photo, etc)

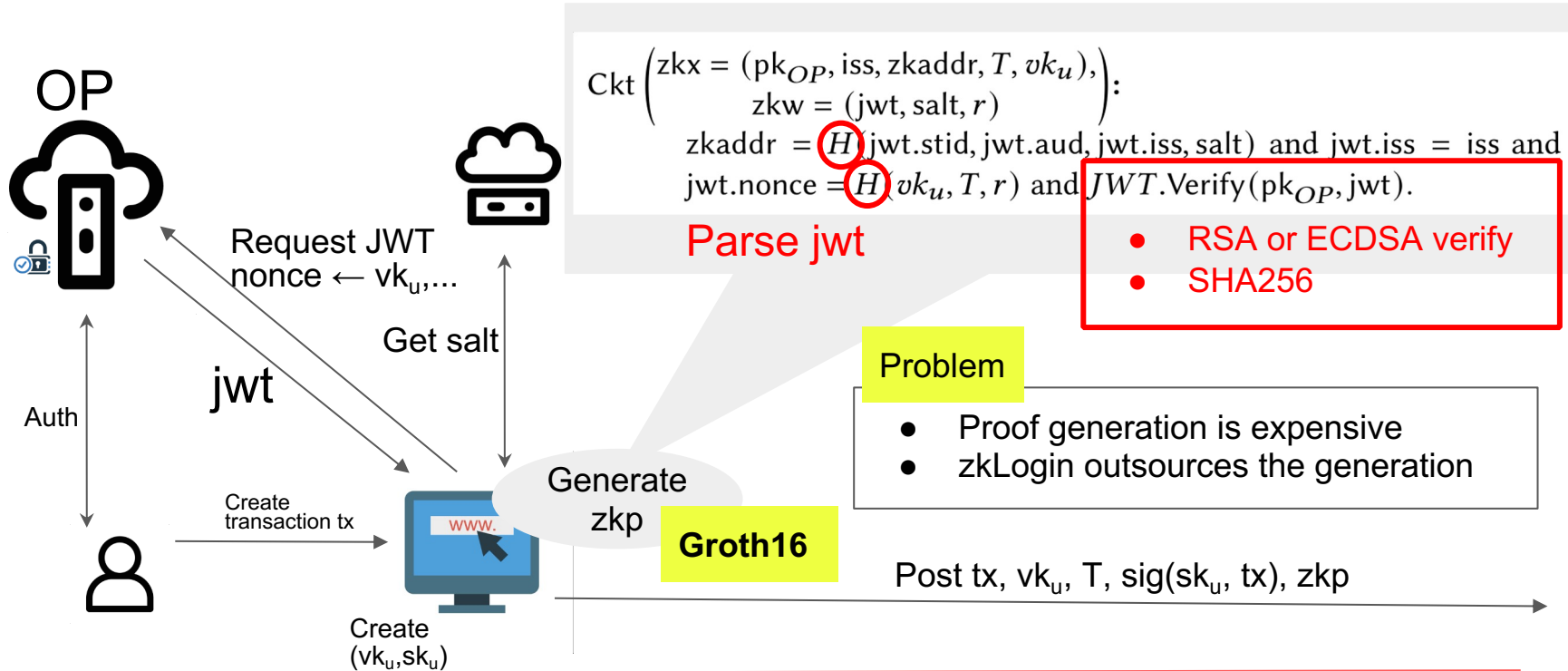
# New solution

## zkLogin: Privacy-Preserving Blockchain Authentication with Existing Credentials

Foteini Baldimtsi	Konstantinos Kryptos Chalkias	Yan Ji	Jonas Lindstrøm
Mysten Labs, George Mason University	Mysten Labs	Cornell University	Mysten Labs
Deepak Maram	Arnab Roy	Mahdi Sedaghat	Joy Wang
Mysten Labs	Mysten Labs	imec-COSIC, KU Leuven	Mysten Labs

@ CCS 2024

# zkLogin flow

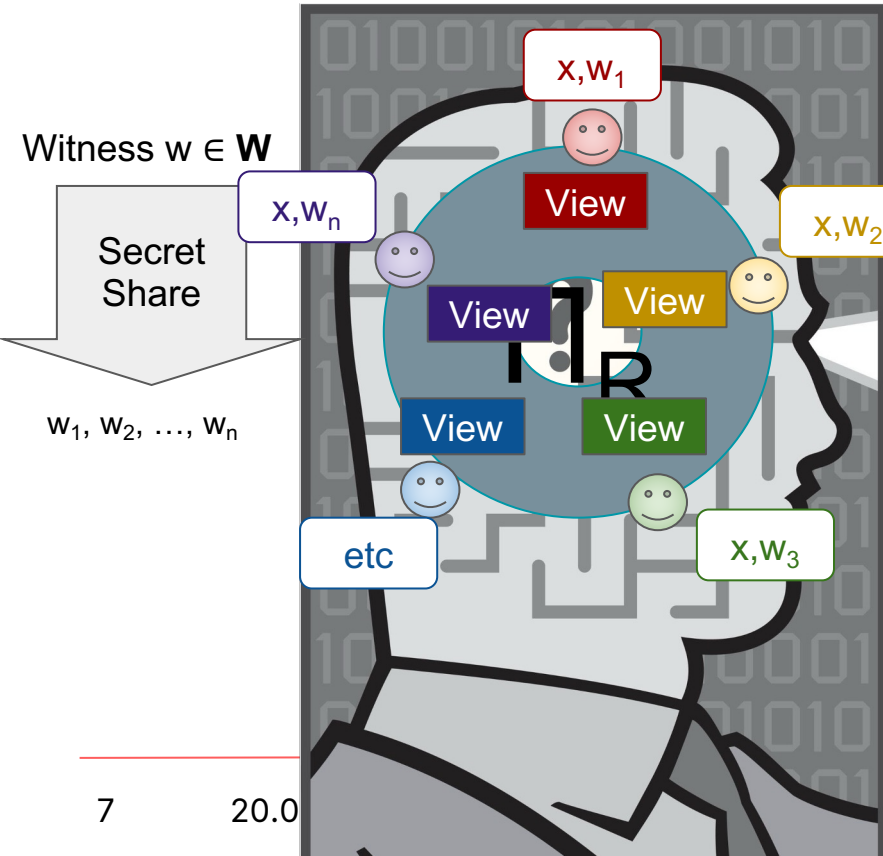


# Motivation + “Research question”

- Alhabill blockchain: <https://github.com/alhabill-org/alhabill>
  - Sharded ledger with heterogeneous rules; fast consensus with “root→shard” trust in validation; programmability via WASM
  - Was developed by Guardtime
- Can the proof be generated in browser?
  - By using some other NIZK technique, perhaps non-succinct?
- CHES: Cyber Security Excellence Hub in Estonia and South Moravia
  - Partners include Cybernetica, Guardtime, [many more]



# MPC in the Head (MPCitH)



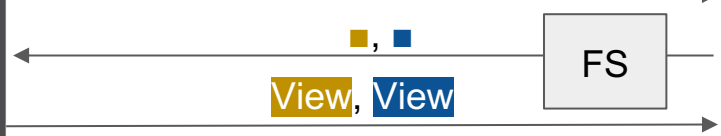
Relation  $R$ , where  $R : \mathbf{X} \times \mathbf{W} \rightarrow \{0,1\}$

Instance  $x \in \mathbf{X}$

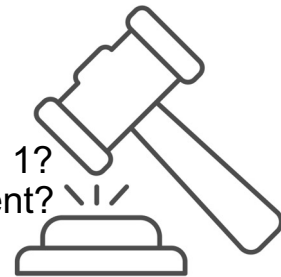
Secure against **two passive** corruptions

$$\prod_R(x; w_1, \dots, w_n) = \prod_{R(x, w)}$$

Com(View), Com(View), ..., Com(View)



- Output = 1?
- Consistent?



Soundness error  $< 1 - 1/n^2$

# Why would MPCitH be suitable?

- Computations in “zkLogin relation” are heterogeneous:
  - Signature verification: large fields
  - Hash function (in signature verification): bits
    - More freedom to choose some other hash functions
  - Parsing: characters, i.e. 8-bit values
  - Base64 encoding: 6- and 8-bit values
- Passively secure MPC can match that heterogeneity

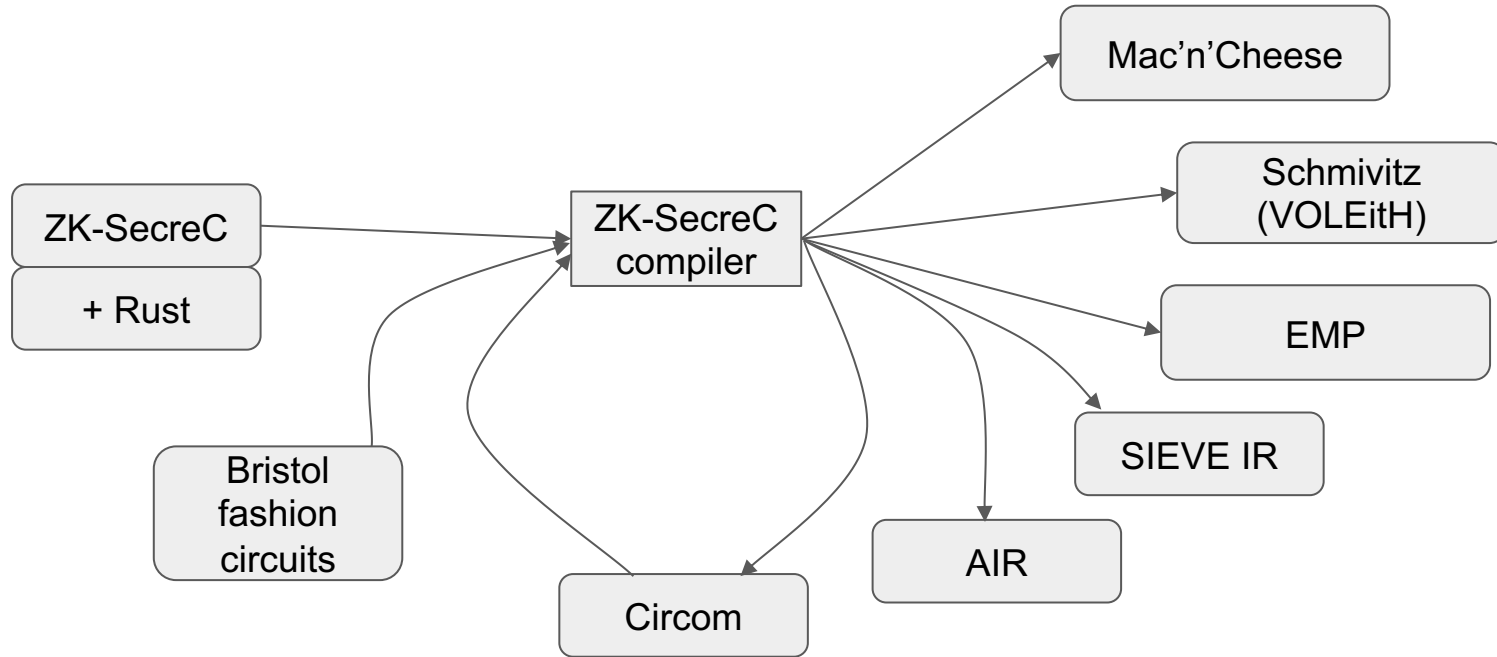
# MPCitH implementations in Github

- ZKB++
  - Very simple protocol
  - A bit of documentation
  - Can read the paper and the code in parallel
    - Consisted of ca. 6k LoC of C++
- Katz-Kolesnikov-Wang
  - No documentation. Could not get it to run

# ZK-SecreC

- DSL for statements one may want to prove in zero-knowledge
- Information-flow type system, with two dimensions:
  - Is a value known only to **prover**, or also to **verifier**? Or is it a **public constant**?
  - Is a value held locally? Or {{managed by the ZKP technique}}?
- Imperative language. Some useful control-flow constructions
- A variety of data types (for “integers”)

# ZK-SecreC toolchain



# Support for many rings in ZKB++

- Supported so far:  $\mathbb{Z}_M$ ,  $M$  is a  $k$ -bit prime, where  $k \in \{3, \dots, 272\}$
- Also supported so far:  $\text{GF}(2^k)$ , where  $k \in \{3, 17, 33, 65\}$

## Our additions

- $\mathbb{Z}_p$  for  $p$  the size of P-256 & P-384 prime fields, and EC groups
- $\mathbb{Z}_M$  for  $M \in \{2^{32}, 2^{64}\}$
- $\mathbb{Z}_2^k$  for  $k \in \{1, \dots, 64, 256, 384\}$  (specialized  $\mathbb{Z}_2^{32}$  and  $\mathbb{Z}_2^{64}$ )
- Elliptic curve *groups* (specialized P-256 & P-384)
- The use of several rings at the same time
  - MPCitH protocol to convert between rings



# Support for rings in ZK-SecreC type system

$n :: \text{Nat}$      $\text{uint} :: \text{Nat} \rightarrow \text{Unqualified}$   
 $\infty :: \text{Nat}$      $\text{bool} :: \text{Nat} \rightarrow \text{Unqualified}$

- $\text{uint}[n]$  is  $\mathbb{Z}_n$ .  $\text{uint}[\infty]$  is  $\mathbb{Z}$
- $\text{bool}[n]$  is booleans, but compatible with  $\text{uint}[n]$

The kind for “pure” data types

$\text{bitwise} :: \text{Nat} \rightarrow \text{Ring}$

$\text{plain} :: \text{Nat} \rightarrow \text{Ring}$

$\text{int} :: \text{Ring} \rightarrow \text{Unqualified}$

$\text{bin} :: \text{Ring} \rightarrow \text{Unqualified}$

- $\text{uint}[n] \equiv \text{int}[\text{plain}[n]]$
- $\text{bool}[n] \equiv \text{bin}[\text{plain}[n]]$

- ZK-SecreC uses SIEVE IR’s “Circuit Configuration Communication” syntax
  - Compiler learns capabilities of the ZKP back-end
    - rings, conversions, plugins...
- This has been extended

Elliptic curve groups: “undocumented API calls”

# Interesting parts of the relation

$$\text{Ckt} \left( \begin{array}{l} \text{zcx} = (\text{pk}_{OP}, \text{iss}, \text{zkaddr}, T, \text{vk}_u), \\ \text{zkw} = (\text{jwt}, \text{salt}, r) \end{array} \right):$$

$\text{zkaddr} = H(\text{jwt.stid}, \text{jwt.aud}, \text{jwt.iss}, \text{salt})$  and  $\text{jwt.iss} = \text{iss}$  and  
 $\text{jwt.nonce} = H(\text{vk}_u, T, r)$  and  $\text{JWT.Verify}(\text{pk}_{OP}, \text{jwt})$ .

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- Parsing JWT (correct JSON only, coming from OP)
  - Slicing an array, running checks on that slice
- Base64 encoding / decoding
- SHA-256 computation
- ECDSA verification (no RSA)

# Our contributions to implementing these parts

## Array slicing

- Recursive approach to pointing out the location of the slice
- Asymptotically decreases the size of the relation

## Base64

- More efficient formula for conversion
- Bit-parallel execution

## ECDSA verification

- Nothing in particular...
  - Leverage the implementation of various rings for ZKB++
  - Sharing over EC group reduces proof size

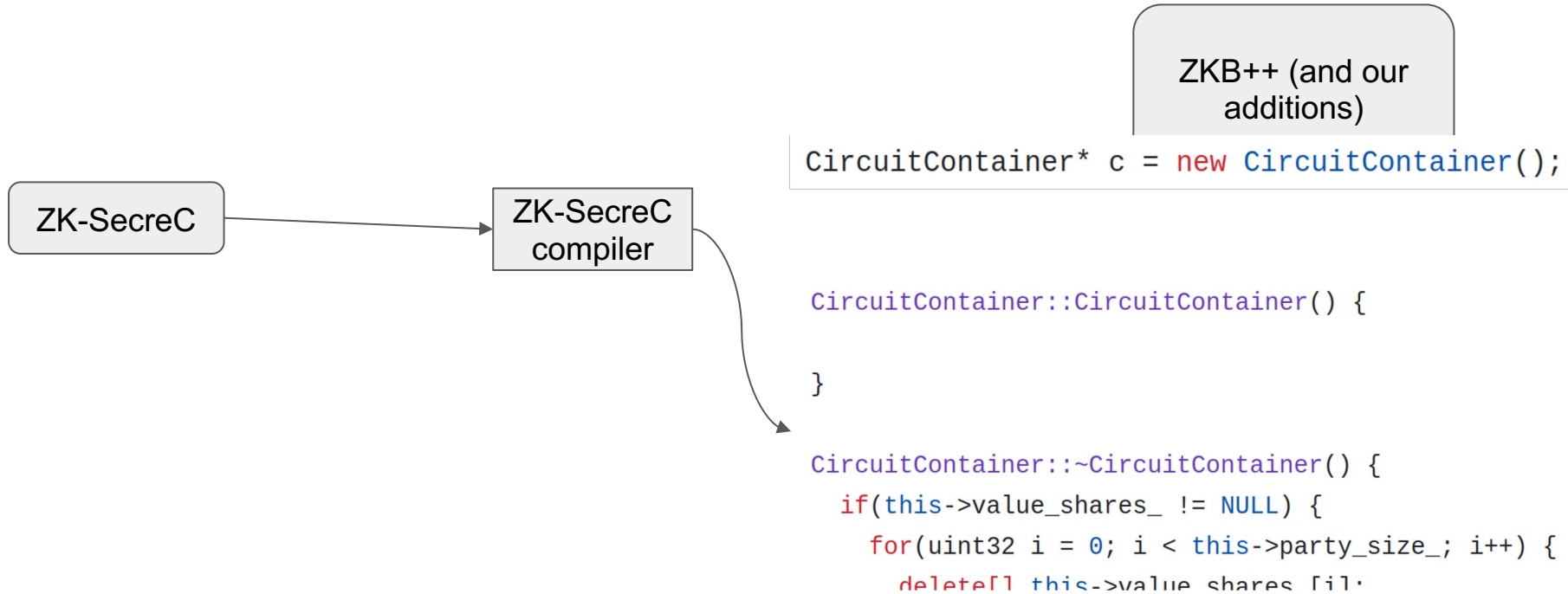
## SHA-256

- Bit-parallel execution of the rounds of the compression function
- Bit-parallel execution of block finishes
  - Careful scheduling, supported by ZK-SecretC

## More additions to ZKB++

- Multiple instance and witness values
- Multiple output values
  - $\Pi_R$  for the relation R “accepts” if all outputs are 0
- Declassification and random number generation in the MPCitH protocol
- Serialization and deserialization of the proof

# ZK-SecreC — ZKB++ integration



# Benchmarks (local execution)

	prover	verifier	witn. expand	size
init + parse + check + base64	0.45s	0.32s	0.09s	15.8 MB
... + SHA	1.49s	1.09s	0.12s	45.5 MB
Full circuit with 256-bit ECDSA in ZKSC	2.06s	1.48s	0.18s	66.2 MB
Full circuit with 256-bit ECDSA in ZKB++	2.08s	1.5s	0.33s	49.1 MB
Full circuit with 384-bit ECDSA in ZKSC	2.54s	1.77s	0.25s	90.1 MB
Full circuit with 384-bit ECDSA in ZKB++	3.01s	2.11s	0.6s	50.8 MB

Execution on a  
2-year old laptop

1472-byte payload. Fixed-length payloads and header. Variable-length adds ca 10%

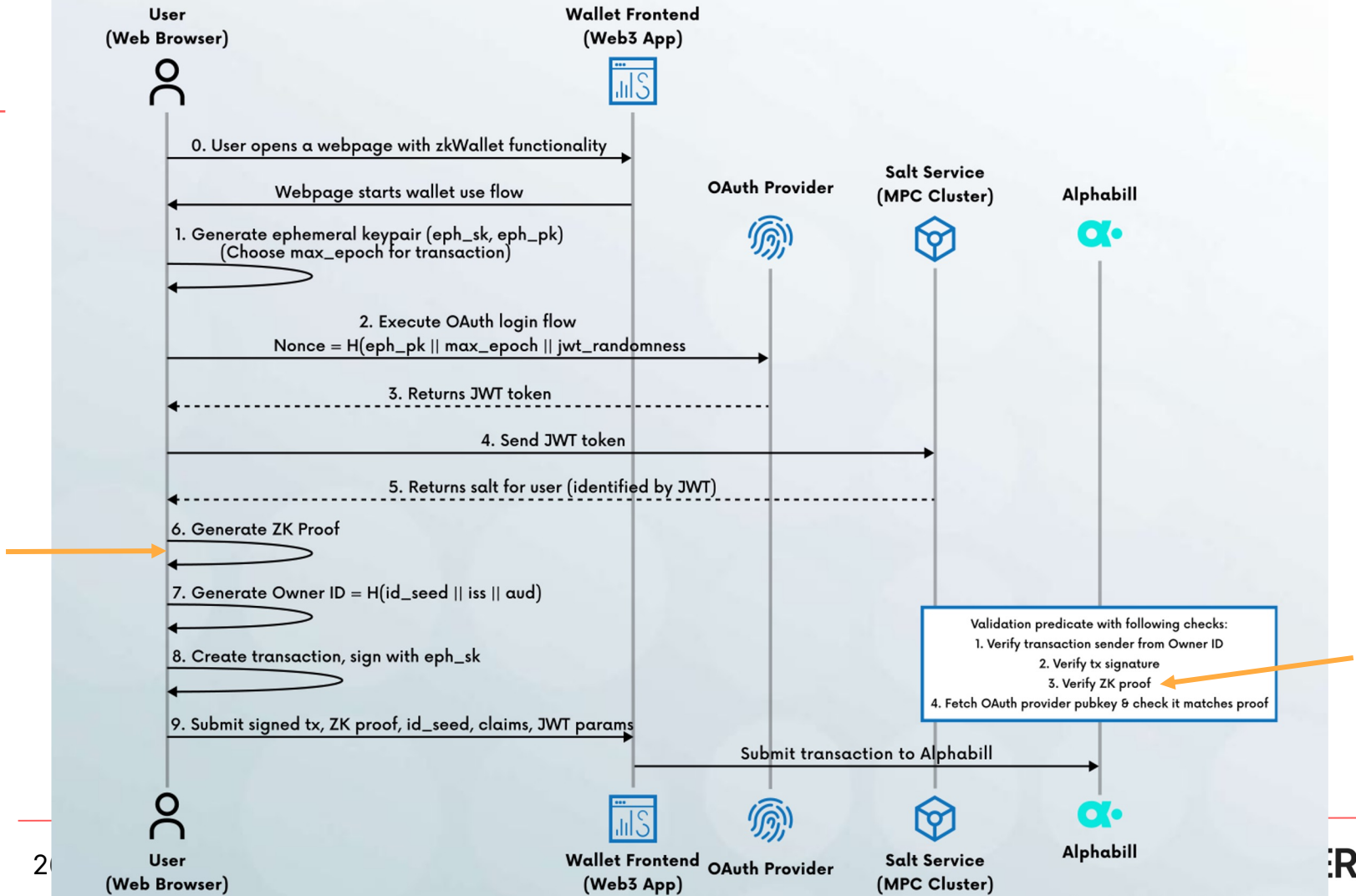
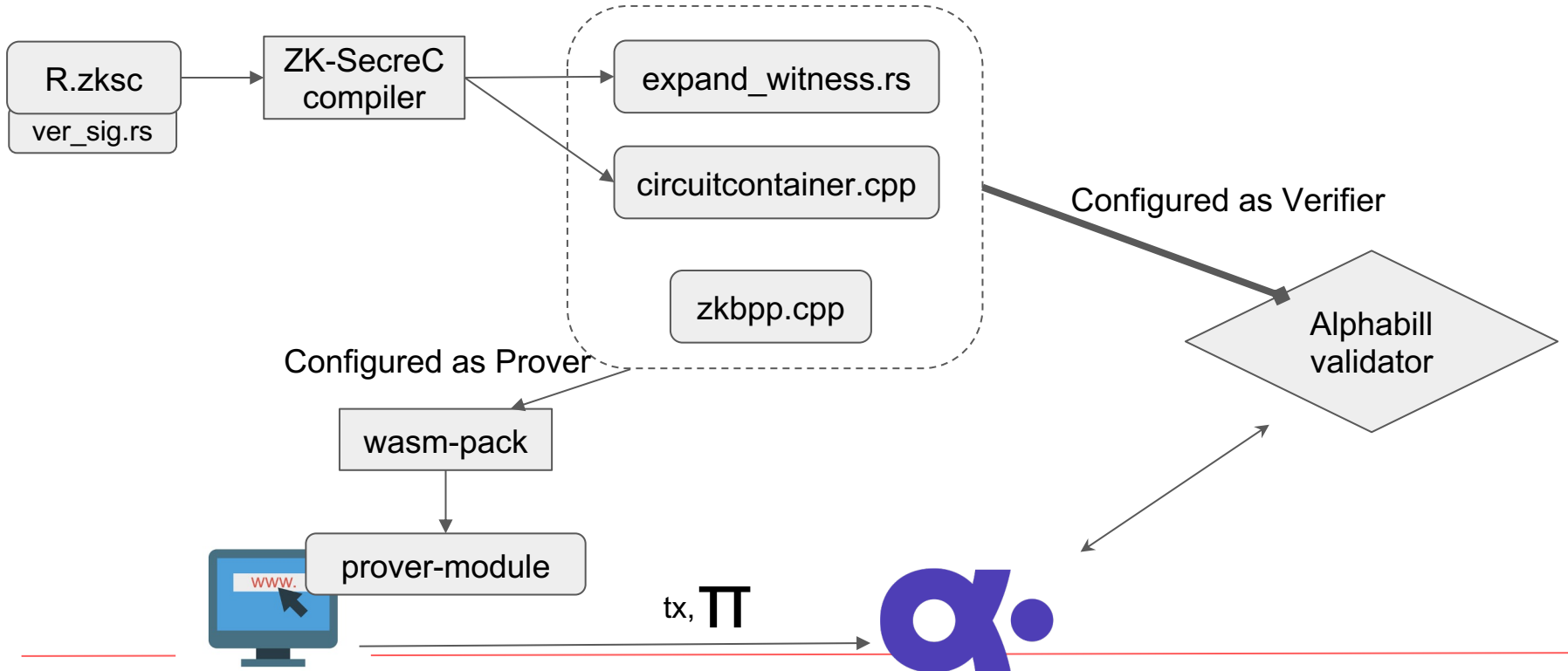


Figure 1. zkLogin workflow overview

# Running in browser + blockchain




# Benchmarks (browser)


Full circuit with 256-bit ECDSA in ZKB++


	witn. exp.	Prover
Firefox 150.0.1 (64bit)	0.4s	4.5s
Chrome 147.0.7727.137 (64-bit)	0.4s	4.3s


CPU: AMD Ryzen 7 PRO 8840U  
OS: NixOS, kernel 7.0.3


# Thank you


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